

Eliminating **letcc** from Labelled Natural Deduction Proofs

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1 Language

Types:

$$A, B ::= P \mid A \wedge B \mid A \vee B \mid A \supset B \mid \top \mid \perp$$

Terms:

$$\begin{aligned} M, N ::= & \langle M_1, M_2 \rangle \mid \pi_i M \mid \mathbf{inj}_i M \\ & \mid (\mathbf{case}_p M \mathbf{of} x_1.M_1 \mid x_2.M_2) \\ & \mid \lambda_a x.M \mid M_1 M_2 \mid \langle \rangle \mid \mathbf{abort}_q M \\ & \mid \mathbf{letcc} u \mathbf{in} M \mid \mathbf{throw} M \mathbf{to} u \end{aligned}$$

Answers:

$$\alpha ::= M \downarrow \mid M \triangleright u \mid (M_p ? x_1.\alpha_1 \mid x_2.\alpha_2) \mid M!_p$$

Pre-answers:

$$\begin{aligned} \beta ::= & \langle \alpha_1, \alpha_2 \rangle \mid \pi_i \alpha \mid \mathbf{inj}_i \alpha \\ & \mid (\mathbf{case}_p \alpha \mathbf{of} x_1.\alpha_1 \mid x_2.\alpha_2) \\ & \mid \lambda_a x.\alpha \mid \alpha_1 \alpha_2 \mid \langle \rangle \mid \mathbf{abort}_q \alpha \\ & \mid \mathbf{letcc} u \mathbf{in} \alpha \mid \mathbf{throw} \alpha \mathbf{to} u \end{aligned}$$

1.1 Refinements of the Syntactic Class of Answers

a -Pure Answers:

$$\alpha^a ::= M \downarrow \mid (M^?_a x_1.\alpha_1^a \mid x_2.\alpha_2^a) \mid M!_a$$

a -Separated Answers: (for $q \not\leq a$)

$$\alpha^{\bar{a}} ::= \alpha^a \mid M \triangleright u \mid (M^?_q x_1.\alpha_1^{\bar{a}} \mid x_2.\alpha_2^{\bar{a}}) \mid M!_q$$

a -Pre-separated Answers:

$$\beta^a ::= (M^?_a x_1.\alpha_1^{\bar{a}} \mid x_2.\alpha_2^{\bar{a}})$$

2 Typing

2.1 Terms and Pre-answers

$$\begin{array}{c}
X ::= M \mid \alpha \\
\\
\frac{x : A[p] \in \Gamma \quad \Gamma \vdash p \leq q}{\Gamma; \Delta \vdash x : A[q]} \\
\\
\frac{\Gamma; \Delta \vdash X_1 : A[p] \quad \Gamma; \Delta \vdash X_2 : B[p]}{\Gamma; \Delta \vdash \langle X_1, X_2 \rangle : A \wedge B[p]} \\
\\
\frac{\Gamma; \Delta \vdash X : A_1 \wedge A_2[p]}{\Gamma; \Delta \vdash \pi_i X : A_i[p]} \\
\\
\frac{\Gamma; \Delta \vdash X : A_i[p]}{\Gamma; \Delta \vdash \mathbf{inj}_i X : A_1 \vee A_2[p]} \\
\\
\frac{\Gamma, x_1 : C_1[q]; \Delta \vdash X_1 : A[p] \quad \Gamma, x_2 : C_2[q]; \Delta \vdash X_2 : A[p]}{\Gamma; \Delta \vdash (\mathbf{case}_q X \mathbf{of} x_1.X_1 \mid x_2.X_2) : A[p]} \\
\\
\frac{\Gamma, x : A[a], a \geq p; \Delta \vdash M : B[a]}{\Gamma; \Delta \vdash \lambda_a x.M : A \rightarrow B[p]} \\
\\
\frac{\Gamma, x : A[a], a \geq p; \Delta \vdash \alpha^{\bar{a}} : B[a]}{\Gamma; \Delta \vdash \lambda_a x.\alpha^{\bar{a}} : A \rightarrow B[p]} \\
\\
\frac{\Gamma; \Delta \vdash X_1 : A \supset B[p] \quad \Gamma; \Delta \vdash X_2 : A[q] \quad \Gamma \vdash q \geq p}{\Gamma; \Delta \vdash X_1 X_2 : B[q]} \\
\\
\frac{}{\Gamma; \Delta \vdash \langle \rangle : \top[p]} \\
\\
\frac{\Gamma; \Delta \vdash X : \perp[q]}{\Gamma; \Delta \vdash \mathbf{abort}_q X : A[p]}
\end{array}$$

2.2 Answers

$$\begin{array}{c}
\frac{\Gamma \downarrow_{\leq p} \vdash M : A}{\Gamma; \Delta \vdash M \downarrow : A[p]} \\
\\
\frac{\Gamma \downarrow_{\leq q} \vdash M : \perp}{\Gamma; \Delta \vdash M !_q : A[p]} \\
\\
\frac{\Gamma \downarrow_{\leq q} \vdash M : B \quad u : B[q] \in \Delta}{\Gamma; \Delta \vdash M \triangleright u : A[p]}
\end{array}$$

$$\frac{\Gamma, x_1 : C_1[q]; \Delta \vdash \alpha_1 : A[p] \quad \Gamma, x_2 : C_2[q]; \Delta \vdash \alpha_2 : A[p]}{\Gamma; \Delta \vdash (M^?_q x_1. \alpha_1 \mid x_2. \alpha_2) : A[p]}$$

3 Termification

$$\boxed{\alpha^a \dashrightarrow_T M}$$

When J_1, J_2 are judgments and R is a relation, the symbol

$$\mathbf{PP} \left\{ \begin{array}{l} J_1(X) \\ X \ R \ Y \\ J_2(Y) \end{array} \right.$$

means that both of the following are true:

- (‘progress’) If $J_1(X)$, then there is a Y such that $X \ R \ Y$.
- (‘preservation’) If $J_1(X)$ and $X \ R \ Y$, then $J_2(Y)$.

For instance, we claim of the termification relation that it satisfies

$$\mathbf{PP} \left\{ \begin{array}{l} \Gamma, \vec{x} : \vec{A}[a], a \geq p; \Delta \vdash \alpha^a : B[a] \\ \alpha_a \dashrightarrow_T M \\ \Gamma \downarrow_{\leq p}, \vec{x} : \vec{A} \vdash M : B \end{array} \right.$$

(assuming that $a \notin \Gamma, \Delta$) when defined by

$$\frac{}{M \downarrow \dashrightarrow_T M}$$

$$\frac{}{M!_a \dashrightarrow_T \mathbf{abort} \ M}$$

$$\frac{\alpha_1^a \dashrightarrow_T M_1 \quad \alpha_2^a \dashrightarrow_T M_2}{M^?_a x_1. \alpha_1^a \mid x_2. \alpha_2^a \dashrightarrow_T \mathbf{case} \ M \ \mathbf{of} \ M_1 \mid M_2}$$

4 Separation

$$\boxed{\beta^a \dashrightarrow_S \alpha^{\bar{a}}}$$

$$\mathbf{PP} \left\{ \begin{array}{l} \Gamma; \Delta \vdash \beta^a : A[p] \\ \beta^a \dashrightarrow_S \alpha^{\bar{a}} \\ \Gamma; \Delta \vdash \alpha^{\bar{a}} : A[p] \end{array} \right.$$

$$\frac{M?_a x_1. \alpha_{11}^{\bar{a}} \mid x_2. \alpha_2^{\bar{a}} \dashrightarrow_S M'_1 \quad M?_a x_1. \alpha_{12}^{\bar{a}} \mid x_2. \alpha_2^{\bar{a}} \dashrightarrow_S M'_2}{(M?_a x_1. (N?_q x_{11}. \alpha_{11}^{\bar{a}} \mid x_{12}. \alpha_{12}^{\bar{a}}) \mid x_2. \alpha_2^{\bar{a}}) \dashrightarrow_S N?_q x_{11}. M'_1 \mid x_{12}. M'_2}$$

$$\frac{M?_a x_1. \alpha_1^{\bar{a}} \mid x_2. \alpha_{21}^{\bar{a}} \dashrightarrow_S M'_1 \quad M?_a x_1. \alpha_1^{\bar{a}} \mid x_2. \alpha_{22}^{\bar{a}} \dashrightarrow_S M'_2}{(M?_a x_1. \alpha_1^{\bar{a}} \mid x_2. (N?_q x_{21}. \alpha_{21}^{\bar{a}} \mid x_{22}. \alpha_{22}^{\bar{a}})) \dashrightarrow_S N?_q x_{21}. M'_1 \mid x_{22}. M'_2}$$

$$\overline{(M?_a x_1. N!_q \mid x_2. \alpha_2^{\bar{a}}) \dashrightarrow_S N!_q} \quad \overline{(M?_a x_1. \alpha_1^{\bar{a}} \mid x_2. N!_q) \dashrightarrow_S N!_q}$$

$$\overline{(M?_a x_1. N \triangleright u \mid x_2. \alpha_2^{\bar{a}}) \dashrightarrow_S N \triangleright u} \quad \overline{(M?_a x_1. \alpha_1^{\bar{a}} \mid x_2. N \triangleright u) \dashrightarrow_S N \triangleright u}$$

$$\overline{(M?_a x_1. \alpha_1^a \mid x_2. \alpha_2^a) \dashrightarrow_S (M?_a x_1. \alpha_1^a \mid x_2. \alpha_2^a)}$$

$$\boxed{\alpha \hookrightarrow_S \alpha^{\bar{a}}}$$

$$\mathbf{PP} \left\{ \begin{array}{l} \Gamma; \Delta \vdash \alpha' : A[p] \\ \alpha' \dashrightarrow_S \alpha^{\bar{a}} \\ \Gamma; \Delta \vdash \alpha^{\bar{a}} : A[p] \end{array} \right.$$

$$\frac{\overline{M \downarrow \hookrightarrow_S M \downarrow} \quad \overline{M!_q \hookrightarrow_S M!_q} \quad \overline{M \triangleright u \hookrightarrow_S M \triangleright u}}{\frac{\alpha'_1 \hookrightarrow_S \alpha_1^{\bar{a}} \quad \alpha'_2 \hookrightarrow_S \alpha_2^{\bar{a}} \quad q \not\leq a}{M?_q x_1. \alpha'_1 \mid x_2. \alpha'_2 \hookrightarrow_S M?_q x_1. \alpha_1^{\bar{a}} \mid x_2. \alpha_2^{\bar{a}}}}{\frac{\alpha'_1 \hookrightarrow_S \alpha_1^{\bar{a}} \quad \alpha'_2 \hookrightarrow_S \alpha_2^{\bar{a}} \quad M?_a x_1. \alpha_1^{\bar{a}} \mid x_2. \alpha_2^{\bar{a}} \dashrightarrow_S \alpha^{\bar{a}}}{M?_a x_1. \alpha'_1 \mid x_2. \alpha'_2 \hookrightarrow_S \alpha^{\bar{a}}}}$$

5 letcc Elimination

$$\boxed{\beta \dashrightarrow \alpha}$$

$$\mathbf{PP} \left\{ \begin{array}{l} \Gamma; \Delta \vdash \beta : A[p] \\ \beta \dashrightarrow \alpha \\ \Gamma; \Delta \vdash \alpha : A[p] \end{array} \right.$$

$$\boxed{M \hookrightarrow \alpha}$$

$$\mathbf{PP} \left\{ \begin{array}{l} \Gamma; \Delta \vdash M : A[p] \\ M \hookrightarrow \alpha \\ \Gamma; \Delta \vdash \alpha : A[p] \end{array} \right.$$

5.1 Conjunction Introduction

$$\frac{M_1 \hookrightarrow \alpha_1 \quad M_2 \hookrightarrow \alpha_2 \quad \langle \alpha_1, \alpha_2 \rangle \dashrightarrow \alpha'}{\langle M_1, M_2 \rangle \hookrightarrow \alpha'}$$

$$\overline{\langle M_1 \downarrow, M_2 \downarrow \rangle \dashrightarrow \langle M_1, M_2 \rangle \downarrow}$$

$$\overline{\langle M!_q, \alpha \rangle \dashrightarrow M!_q} \quad \overline{\langle \alpha, M!_q \rangle \dashrightarrow M!_q}$$

$$\overline{\langle M \triangleright u, \alpha \rangle \dashrightarrow M \triangleright u} \quad \overline{\langle \alpha, M \triangleright u \rangle \dashrightarrow M \triangleright u}$$

$$\overline{\langle \alpha_1, \alpha \rangle \dashrightarrow \alpha'_1} \quad \overline{\langle \alpha_2, \alpha \rangle \dashrightarrow \alpha'_2}$$

$$\overline{\langle (M?_q x_1. \alpha_1 \mid x_2. \alpha_2), \alpha \rangle \dashrightarrow M?_q x_1. \alpha'_1 \mid x_2. \alpha'_2}$$

$$\overline{\langle \alpha, \alpha_1 \rangle \dashrightarrow \alpha'_1} \quad \overline{\langle \alpha, \alpha_2 \rangle \dashrightarrow \alpha'_2}$$

$$\overline{\langle \alpha, (M?_q x_1. \alpha_1 \mid x_2. \alpha_2) \rangle \dashrightarrow M?_q x_1. \alpha'_1 \mid x_2. \alpha'_2}$$

5.2 Conjunction Elimination

$$\frac{M \hookrightarrow \alpha \quad \pi_i \alpha \dashrightarrow \alpha'}{\pi_i M \hookrightarrow \alpha'}$$

$$\overline{\pi_i(M \downarrow) \dashrightarrow (\pi_i M) \downarrow} \quad \overline{\pi_i(M!_q) \dashrightarrow M!_q} \quad \overline{\pi_i(M \triangleright u) \dashrightarrow M \triangleright u}$$

$$\overline{\pi_i \alpha_1 \dashrightarrow \alpha'_1} \quad \overline{\pi_i \alpha_2 \dashrightarrow \alpha'_2}$$

$$\overline{\pi_i(M?_q x_1. \alpha_1 \mid x_2. \alpha_2) \dashrightarrow M?_q x_1. \alpha'_1 \mid x_2. \alpha'_2}$$

5.3 Disjunction Introduction

$$\begin{array}{c}
\frac{M \hookrightarrow \alpha \quad \mathbf{inj}_i \alpha \dashrightarrow \alpha'}{\mathbf{inj}_i M \hookrightarrow \alpha'} \\
\\
\frac{}{\mathbf{inj}_i(M \downarrow) \dashrightarrow (\mathbf{inj}_i M) \downarrow} \quad \frac{}{\mathbf{inj}_i(M !_q) \dashrightarrow M !_q} \quad \frac{}{\mathbf{inj}_i(M \triangleright u) \dashrightarrow M \triangleright u} \\
\\
\frac{\mathbf{inj}_i \alpha_1 \dashrightarrow \alpha'_1 \quad \mathbf{inj}_i \alpha_2 \dashrightarrow \alpha'_2}{\mathbf{inj}_i(M ?_q x_1. \alpha_1 \mid x_2. \alpha_2) \dashrightarrow M ?_q x_1. \alpha'_1 \mid x_2. \alpha'_2}
\end{array}$$

5.4 Disjunction Elimination

$$\begin{array}{c}
\frac{M \hookrightarrow \alpha \quad M_1 \hookrightarrow \alpha_1 \quad M_2 \hookrightarrow \alpha_2 \quad \mathbf{case}_q \alpha \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow \alpha'}{\mathbf{case}_q M \text{ of } x_1. M_1 \mid x_2. M_2 \hookrightarrow \alpha'} \\
\\
\frac{}{\mathbf{case}_q M ! \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow M !} \\
\\
\frac{}{\mathbf{case}_q M \triangleright u \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow M \triangleright u} \\
\\
\frac{\mathbf{case}_q \alpha'_1 \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow \alpha''_1 \quad \mathbf{case}_q \alpha'_2 \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow \alpha''_2}{\mathbf{case}_q (M ?_r y_1. \alpha'_1 \mid y_2. \alpha'_2) \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow M ?_r y_1. \alpha''_1 \mid y_2. \alpha''_2} \\
\\
\frac{}{\mathbf{case}_q M \downarrow \text{ of } x_1. \alpha_1 \mid x_2. \alpha_2 \dashrightarrow M ?_q x_1. \alpha_1 \mid x_2. \alpha_2}
\end{array}$$

5.5 Implication Introduction

$$\begin{array}{c}
\frac{M \hookrightarrow \alpha' \quad \alpha' \hookrightarrow_S \alpha^{\bar{a}} \quad \lambda_a x. \alpha^{\bar{a}} \dashrightarrow \alpha''}{\lambda_a x. M \hookrightarrow \alpha''} \\
\\
\frac{\alpha^a \dashrightarrow_T M}{\lambda_a x. \alpha^a \dashrightarrow \lambda x. M \downarrow} \quad \frac{}{\lambda_a x. M !_q \dashrightarrow M !_q} \quad \frac{}{\lambda_a x. M \triangleright u \dashrightarrow M \triangleright u} \\
\\
\frac{\lambda_a x. \alpha^{\bar{a}}_1 \dashrightarrow \alpha'_1 \quad \lambda_a x. \alpha^{\bar{a}}_2 \dashrightarrow \alpha'_2}{\lambda_a x. M ?_q x_1. \alpha^{\bar{a}}_1 \mid x_2. \alpha^{\bar{a}}_2 \dashrightarrow M ?_q x_1. \alpha'_1 \mid x_2. \alpha'_2}
\end{array}$$

5.6 Implication Elimination

$$\begin{array}{c}
\frac{M_1 \hookrightarrow \alpha_1 \quad M_2 \hookrightarrow \alpha_2 \quad \alpha_1 \ \alpha_2 \dashrightarrow \alpha'}{M_1 \ M_2 \hookrightarrow \alpha'} \\
\\
\frac{}{M_1 \downarrow \ M_2 \downarrow \dashrightarrow (M_1 \ M_2) \downarrow} \\
\\
\frac{}{M!_q \ \alpha \dashrightarrow M!_q} \quad \frac{}{\alpha \ M!_q \dashrightarrow M!_q} \\
\\
\frac{}{(M \triangleright u) \ \alpha \dashrightarrow M \triangleright u} \quad \frac{}{\alpha \ (M \triangleright u) \dashrightarrow M \triangleright u} \\
\\
\frac{\alpha_1 \ \alpha \dashrightarrow \alpha'_1 \quad \alpha_2 \ \alpha \dashrightarrow \alpha'_2}{(M?_q x_1. \alpha_1 \mid x_2. \alpha_2) \ \alpha \dashrightarrow M?_q x_1. \alpha'_1 \mid x_2. \alpha'_2} \\
\\
\frac{\alpha \ \alpha_1 \dashrightarrow \alpha'_1 \quad \alpha \ \alpha_2 \dashrightarrow \alpha'_2}{\alpha \ (M?_q x_1. \alpha_1 \mid x_2. \alpha_2) \dashrightarrow M?_q x_1. \alpha'_1 \mid x_2. \alpha'_2}
\end{array}$$

5.7 Truth Introduction

$$\begin{array}{c}
\frac{\langle \rangle \dashrightarrow \alpha'}{\langle \rangle \hookrightarrow \alpha'} \\
\\
\frac{}{\langle \rangle \dashrightarrow \langle \rangle \downarrow}
\end{array}$$

5.8 Falsehood Elimination

$$\begin{array}{c}
\frac{M \hookrightarrow \alpha \quad \mathbf{abort}_q \ \alpha \dashrightarrow \alpha'}{\mathbf{abort}_q \ M \hookrightarrow \alpha'} \\
\\
\frac{}{\mathbf{abort}_q \ M! \dashrightarrow M!} \\
\\
\frac{}{\mathbf{abort}_q (M \triangleright u) \dashrightarrow M \triangleright u} \\
\\
\frac{\mathbf{abort}_q \ \alpha'_1 \dashrightarrow \alpha''_1 \quad \mathbf{abort}_q \ \alpha'_2 \dashrightarrow \alpha''_2}{\mathbf{abort}_q (M?_r y_1. \alpha'_1 \mid y_2. \alpha'_2) \dashrightarrow M?_r y_1. \alpha''_1 \mid y_2. \alpha''_2} \\
\\
\frac{}{\mathbf{abort}_q (M \downarrow) \dashrightarrow M!_q}
\end{array}$$

5.9 Throw

$$\begin{array}{c}
\frac{M \hookrightarrow \alpha \quad \mathbf{throw} \ \alpha \ \mathbf{to} \ u \dashrightarrow \alpha'}{\mathbf{throw} \ M \ \mathbf{to} \ u \hookrightarrow \alpha'} \\
\\
\frac{}{\mathbf{throw} \ M! \ \mathbf{to} \ u \dashrightarrow M!} \\
\\
\frac{}{\mathbf{throw} \ (M \triangleright v) \ \mathbf{to} \ u \dashrightarrow M \triangleright v} \\
\\
\frac{\mathbf{throw} \ \alpha'_1 \ \mathbf{to} \ u \dashrightarrow \alpha''_1 \quad \mathbf{throw} \ \alpha'_2 \ \mathbf{to} \ u \dashrightarrow \alpha''_2}{\mathbf{throw} \ (M?_{ry_1}.\alpha'_1 \mid y_2.\alpha'_2) \ \mathbf{to} \ u \dashrightarrow M?_{ry_1}.\alpha''_1 \mid y_2.\alpha''_2} \\
\\
\frac{}{\mathbf{throw} \ (M \downarrow) \ \mathbf{to} \ u \dashrightarrow M \triangleright u}
\end{array}$$

5.10 Letcc

$$\begin{array}{c}
\frac{M \hookrightarrow \alpha \quad \mathbf{letcc} \ u \ \mathbf{in} \ \alpha \dashrightarrow \alpha'}{\mathbf{letcc} \ u \ \mathbf{in} \ M \hookrightarrow \alpha'} \\
\\
\frac{}{\mathbf{letcc} \ u \ \mathbf{in} \ M! \dashrightarrow M!} \\
\\
\frac{u \neq v}{\mathbf{letcc} \ u \ \mathbf{in} \ M \triangleright v \dashrightarrow M \triangleright v} \quad \frac{}{\mathbf{letcc} \ u \ \mathbf{in} \ M \triangleright u \dashrightarrow M \downarrow} \\
\\
\frac{\mathbf{letcc} \ u \ \mathbf{in} \ \alpha'_1 \dashrightarrow \alpha''_1 \quad \mathbf{letcc} \ u \ \mathbf{in} \ \alpha'_2 \dashrightarrow \alpha''_2}{\mathbf{letcc} \ u \ \mathbf{in} \ M?_{ry_1}.\alpha'_1 \mid y_2.\alpha'_2 \dashrightarrow M?_{ry_1}.\alpha''_1 \mid y_2.\alpha''_2} \\
\\
\frac{}{\mathbf{letcc} \ u \ \mathbf{in} \ M \downarrow \dashrightarrow M \downarrow}
\end{array}$$